Infinispon

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Who am I?

- Core developer for Infinispan and JBoss Cache
- Contributor and committer on JBoss AS, JGroups, Hibernate, JBoss Portal, etc.



Agenda

- What is Infinispan?
- Relationship with JBoss Cache
- New features

74

• Demo



What is Infinispan?

Highly scalable data grid platform

- 100% open source licensed (LGPL)
- Based on some JBoss Cache code

• JBoss Cache = Tree structured cache

- Replicated using JGroups
- Supports JTA, evictions, cache stores, etc.
- New JSR-107(JCACHE) compatible API
 - Cache extends Map
 - Tree adapter API available for legacy apps

More scalable than JBoss Cache

Internal structures more memory efficient

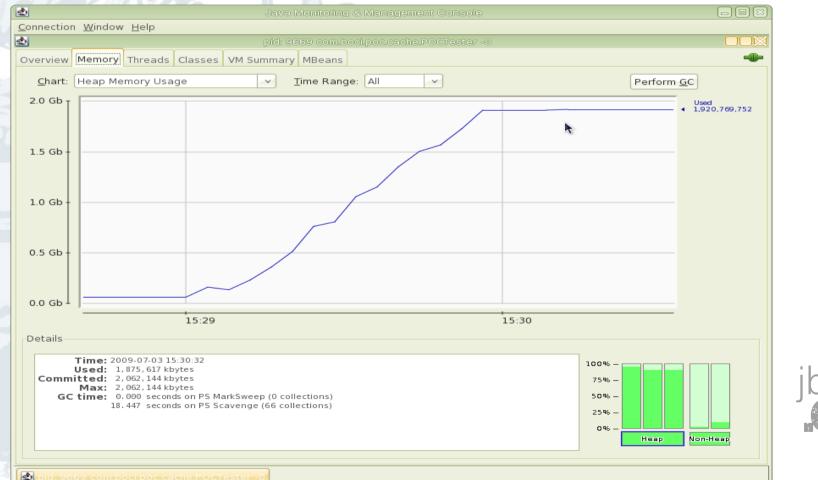
- Tree --> Flat concurrent map
- Eviction queue --> Ordered container
- Marshalling based on JBoss Marshalling
 - Smaller payloads + Poolable streams
- Early benchmarks
 - Significant performance improvements



Memory Consumption Comparison

Test: Put 2 million serial objs into cache

• JBoss Cache 3.1: With 2gb, 1 million objs



Memory Consumption Comparison (2)

Infinispan 4.0.0.Alpha4: With 700mb, 2 million objs

٤.	Java Monitoring & Management Console	
Connection <u>W</u> indow <u>H</u> elp		
<u> </u>	pid: 10125 com.boci.poc.cache.POCTester -0	
Overview Memory Threads Class	es VM Summary MBeans	
<u>C</u> hart: Heap Memory Usage	✓ <u>T</u> ime Range: All ✓	Perform <u>G</u> C
800 Mb T		Used
700 Mb -		4 767,480,2
600 Mb		
500 Mb -		
400 Mb -		
300 мb -	/	
200 Mb -		
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0.0 мb		N
15:34		
Details		
Time: 2009-07-03 15:34:56		100%
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Max: 2,019,520 kbytes		50% -
GC time: 0,000 seconds on PS		
GC time: 0.000 seconds on PS	S Scavenge (14 collections)	25% -
GC time: 0.000 seconds on PS	5 Scavenge (14 collections)	25% -
GC time: 0.000 seconds on PS	S Scavenge (14 collections)	25%

...borrowing best bits from JBoss Cache

- Multiversion Concurrency Control (MVCC)
 - New locking strategy in JBoss Cache 3.0
 - Readers never locked!
 - Writers work on copy of cache entry
- Non-blocking state transfer
 - Senders generate state without stopping
 - Crucial when state is large



New features - distributed cache

- Consistent hash based distribution
 - Will allow us to scale to bigger clusters
- Lightweight, L1 cache for efficient reads
 - On writes, L1 invalidated
- Dynamic rebalancing
- Pluggable consistent hashing algorithms
- Already available in 4.0.0.Alpha5!



New features - asynchronous API

- putAsync(), putIfAbsentAsync()
 - Do not block, return a j.u.c.Future
 - Future.get() blocks till call completes
- Best of both sync and async worlds
 - Future.get() provides sync guarantees
 - Greater parallelism
- Already available in 4.0.0.Alpha5!



New features - Eager locking

By default: locks acquired at commit time

- Problematic if updating a shared counter
- New: Acquiring locks eagerly in cluster
 - Explicit: via API
 - cache.lock(k) // acquire cluster wide lock on k
 Implicit: via configuration
 - Each modification implicitly acquires cluster wide lock if not already held.

Already available in 4.0.0.Alpha5!



New features - client/server module

- Server module = cache wrapper over TCP
- Client module = cache proxy
- Highly pluggable!
 - Transport: XNIO, Netty, etc.
 - Protocols: memcached, custom, etc.
- Failover and load balancing
- Usable with current memcached clients
- Drop-in replacement memcached servers

New features - fine-grained model

- Successor to POJO Cache
- JPA-like interface: persist, find, remove...
- Will not rely on AOP, javassist...etc
 - More robust and easier to use/debug



New features - Others

Query module

- Execute Lucene queries against cache
- Based on JBoss Cache Searchable
- Distributed executors
 - Runnable/Callable executed on data set
 - Moves code, not data, around cluster



The End

- Demo
- Questions?

7.4

